Bounded ultraimaginary independence

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Something for nothing: Independence in arbitrary theories

In tame contexts: Independence notion \Rightarrow Generic sequences

• Stable and simple: Non-forking \Rightarrow Morley sequences

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If so, we can build \bigcup^* -Morley sequences: $(b_i)_{i < \omega}$ s.t. $b_i \bigcup_{\Delta}^* b_{< i}$.

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- Usually want a *total* \downarrow^* -Morley sequence: $(b_i)_{i < \omega}$ s.t. if $I + J \equiv_A^{\text{EM}} b_{<\omega}$, then $J \downarrow_A^* I$.

Q2 Can we build total \downarrow^* -Morley sequences?

Weakest 'reasonable' independence relation:

$$b \stackrel{\mathsf{l}}{\underset{A}{\downarrow}^{\mathsf{a}}} c \Leftrightarrow \operatorname{acl}(Ab) \cap \operatorname{acl}(Ac) = \operatorname{acl}(A)$$

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Weakest 'reasonable' independence relation:

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- Bad news: \bigcup^{a} doesn't seem to mean much in arbitrary theories.

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With elimination of hyperimaginaries we can replace \bigcup^{b} with \bigcup^{a} .

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 Good news: ^{bu} definitely means something.

What something does U^{bu} mean?

 $\operatorname{Autf}(\mathbb{M}/A)$ is the group generated by

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 $b \equiv^{\mathsf{L}}_{\mathsf{A}} b' \text{ iff } b' \in \operatorname{Autf}(\mathbb{M}/\mathsf{A}) \cdot b.$

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TFAE:

$$\bullet \bigcup_{A}^{\mathsf{bu}} c$$

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- $\operatorname{Autf}(\mathbb{M}/A)$ is generated by $\operatorname{Autf}(\mathbb{M}/Ab) \cup \operatorname{Autf}(\mathbb{M}/Ac)$.
- (*Walking*) For any $b' \equiv^{\mathsf{L}}_{A} b$, we have the configuration

$$b_0 \equiv^{\mathsf{L}}_{Ac_1} b_2 \equiv^{\mathsf{L}}_{Ac_3} b_4 \equiv^{\mathsf{L}}_{Ac_5} \cdots b_{n-2} \equiv^{\mathsf{L}}_{Ac_{n-1}} b_n$$

$$c_1 \equiv^{\mathsf{L}}_{Ab_2} c_3 \equiv^{\mathsf{L}}_{Ab_4} c_5 \equiv^{\mathsf{L}}_{Ab_6} \cdots c_{n-1}$$

where $b_0 = b$, $c_1 = c$, and $b_n = b'$.

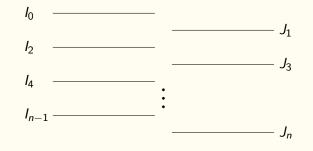
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- Shelah's definition in early simplicity theory: *I* is based on *A* if $I \equiv_A J \Leftrightarrow I \approx_A J$.

What are total U^{bu}-Morley sequences? II

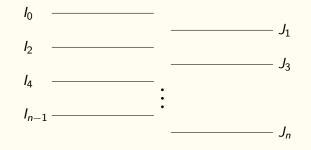
Canonical witnessing configuration: $I \approx_A J$ if and only if we have



where $I_0 = I$, $J_n = J$, and $I_i + J_{i+1}$ and $I_{i+2} + J_{i+1}$ are A-indiscernible.

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Theorem (H.)

 $(b_i)_{i < \omega}$ is a total \downarrow^{bu} -Morley sequence over A iff it is based on $bdd^u(A)$ (i.e. $I \equiv^L_A b_{<\omega} \Leftrightarrow I \approx_A b_{<\omega}$).

Note:
$$I \equiv_{bdd^{u}(A)} J$$
 iff $I \equiv^{L}_{A} J$.

The two questions

Theorem (H.) Yes.

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Proof.

Horrible indiscernible tree combinatorics à la Kaplan-Ramsey.



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There is a 'chain condition': If $(b_i)_{i < \omega}$ is a \bigcup_{a}^{bu} -Morley sequence over A that is Ac-indiscernible, then $c \bigcup_{a}^{bu} b_0$.

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Corollary (H.) $\downarrow^{d} \Rightarrow \downarrow^{bu}$

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Suppose $c extstyle _A^d b$. Find a $extstyle _b^{\text{bu}}$ -Morley sequence $b_{<\omega}$ over A with $b_0 = b$. Since $c extstyle _A^d b$, we may assume that $b_{<\omega}$ is Ac-indiscernible.

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Suppose $c extstyle{bergen}^{d} b$. Find a $extstyle{burgen}^{burgen}$ -Morley sequence $b_{<\omega}$ over A with $b_0 = b$. Since $c extstyle{bergen}^{d} b$, we may assume that $b_{<\omega}$ is Ac-indiscernible. By the chain condition, $c extstyle{burgen}^{burgen} b$.

Corollary of Corollary

In a simple theory, $(b_i)_{i < \omega}$ is a Morley sequence over A if and only if it is a total \bigcup^{bu} -Morley sequence over A.

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In NSOP₁ theories

What about $NSOP_1$ theories?

What about NTP_1 theories?

What about NTP₁ theories?

Proposition (H.)

 $(T \text{ NTP}_1)$ If *I* is a tree Morley sequence over $M \models T$, then *I* is a total $\bigcup_{i=1}^{bu}$ -Morley sequence over *M*.

What about NTP₁ theories?

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Converse?

■ Odd observation: In stable theories, you get a '~_A-distance' of 2. In simple theories, you get 3. And in NTP₁ theories, you get 4.

Q2: Total \bigcup^{bu} -Morley sequences?

Given A and b, can we find a total \bigcup^{bu} -Morley sequence $(b_i)_{i < \omega}$ over A with $b_0 = b$?



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- Does this actually need large cardinals?
- Without any set theoretic hypotheses, we can get a sequence $(b_i)_{i < \omega}$ such that $b_{<i} \bigcup_{A}^{b_u} b_{\geq i}$ for each $i < \omega$.

Applications

Strong witnesses of Lascar strong type

Fix A and b and suppose there is a total $\bigcup_{i=1}^{bu}$ -Morley sequence $I \ni b$. For any b' with $b' \equiv^{L}_{A} b$, we have the configuration b J_1 b Jz I₄ J_{n-1} In h' with $I_0 = I$, $b' \in I_n$, and $I_i + J_{i+1}$ and $I_{i+2} + J_{i+1}$ A-indiscernible for all *i*.

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with $I_0 = I$, $b' \in I_n$, and $I_i + J_{i+1}$ and $I_{i+2} + J_{i+1}$ A-indiscernible for all *i*.

This is similar to a configuration in the proof of the independence theorem.

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Theorems (Shelah, Hrushovski, Kim–Pillay, Ben Yaacov–Chernikov, Kaplan–Ramsey, Simon, Dobrowolski–Kim–Ramsey, etc.)

(*T* nice, maybe) Let $\Sigma(x)$ be an *A*-invariant partial type satisfying a chain condition. Assume that $c \models \Sigma \upharpoonright Aab$ and $b \equiv^{\mathsf{L}}_{A} b'$ and that *a*, *b*, and *b'* are sufficiently independent of one another. Then there exists a $c' \models \Sigma \upharpoonright Aab'$ such that $ac' \equiv_A ac$ and $b'c' \equiv_A bc$.

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 $\Sigma(x)$ is often a *generically prime* filter: If $(b_i)_{i < \omega}$ is *A*-indiscernible and $\Sigma(x) \vdash \varphi(x, b_0) \lor \varphi(x, b_1)$, then $\Sigma(x) \vdash \varphi(x, b_0)$.

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Let $\Sigma(x)$ be A-invariant and generically prime over A. For any a, I, I', and c, if

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 are total \bigcup^{bu} -Morley sequences over A ,

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$$c \models \Sigma \upharpoonright Aab$$
 for all $b \in I$, and

$$|I|, |I'| > 2^{|Aabc| + |T|},$$

then there are $b \in I$, $b' \in I'$, and $c' \models \Sigma \upharpoonright Aab'$ such that $ac' \equiv_A ac$ and $b'c' \equiv_A bc$.

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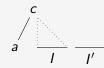
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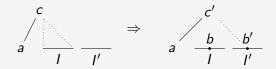
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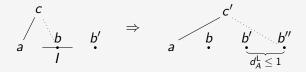
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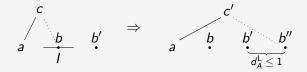


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Can we weaken the generic primality requirement?

Thank you